

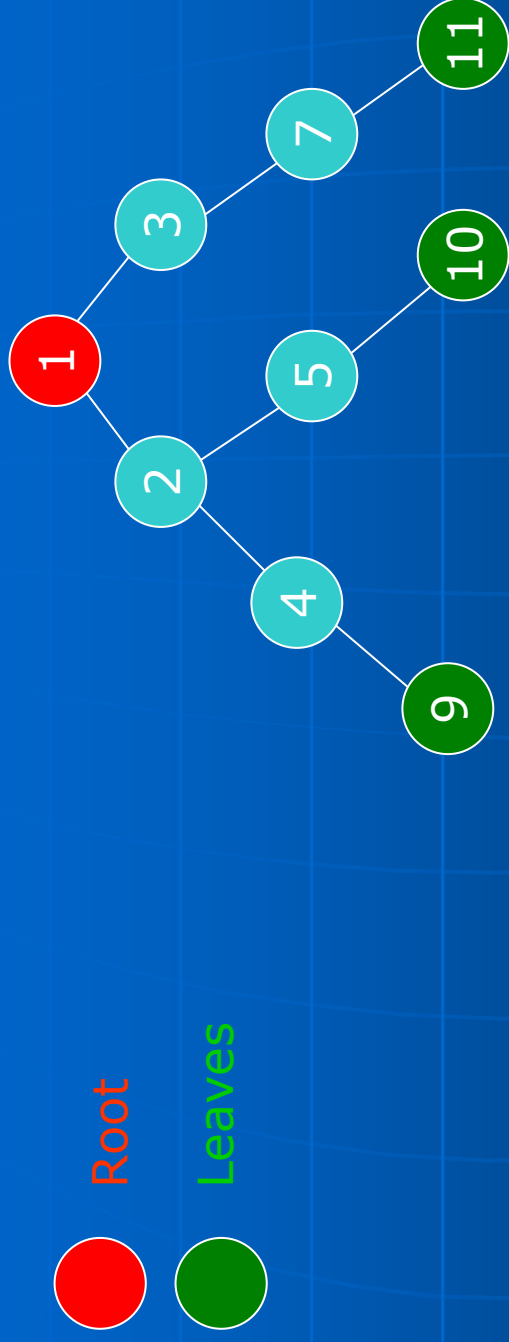
# Trees

By Charl du Plessis

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- Basic Terminology
- Binary Search Trees
- Interval Trees
- Binary Indexed Trees

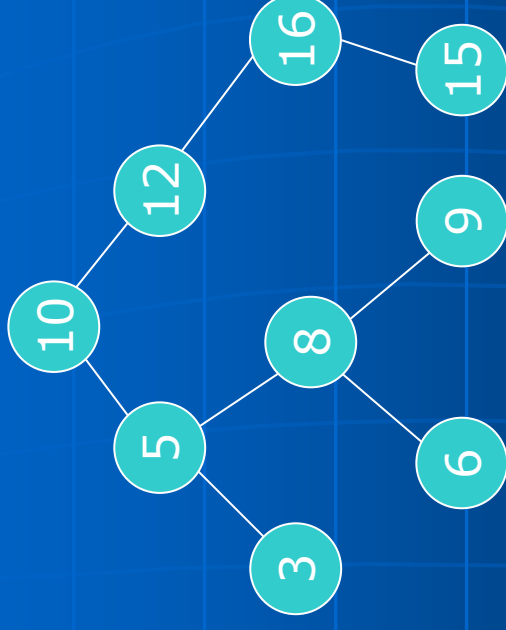
# Basic Terminology



- Connected graph with no cycles
- Unique path from every vertex to every other vertex
- $E = V - 1$  in any tree, where  $E$  and  $V$  are the number of edges and vertices respectively.

# Binary Search Trees

- Structure?
- Operations:  $O(h)$ 
  - query
  - delete
  - insert



Cons:

Degenerate trees make all operations Worst case  $O(n)$   
But can be “self-balanced” such as in  
Red-Black Trees.

C++ `stl <set>` uses self-balancing binary trees

# Interval Trees: Problem

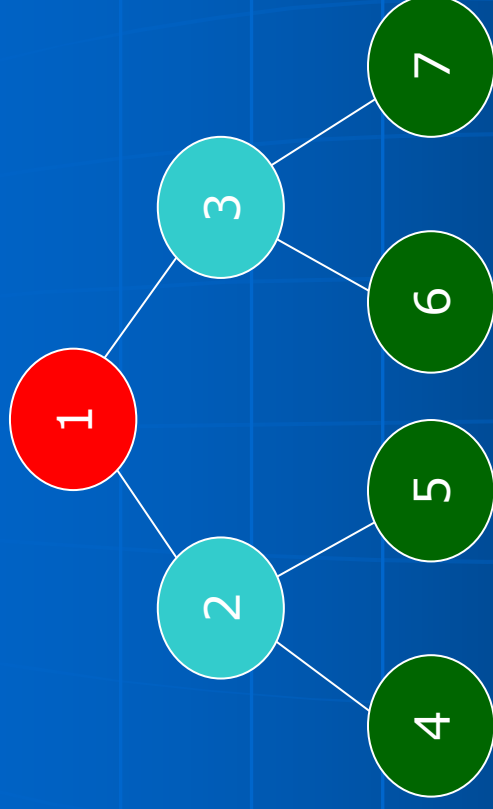
Consider the problem:

- You have a set of  $n$  items;  $1, \dots, n$
- There is some number of each item
- You want to increment the number of each item in range  $[x, y]$  by one. (for each  $i$ ,  $x \leq i \leq y$ )
- And query how many of item  $i$  there are

# Interval Trees: Structure

- Use a binary tree, where each node holds a count over range  $[A, B]$
- Root holds count over  $[1, n]$
- Node ranging over  $[A, B]$  will have children  $[A, (A+B)/2]$  and  $[(A+B)/2+1, B]$
- The leaves of the tree will store count over a single item range

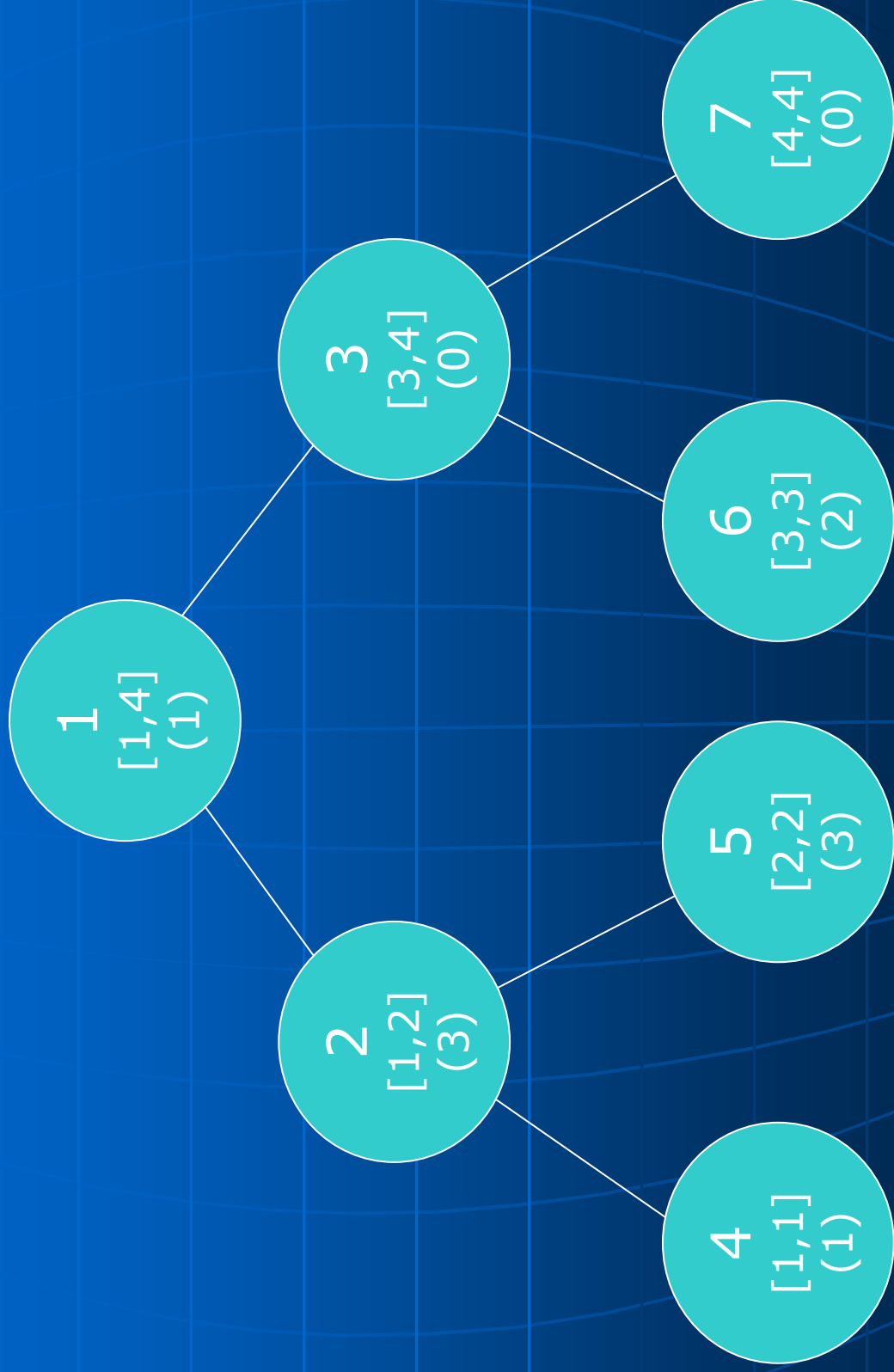
# Interval Trees: Representation



Representation of tree:

- Use an array
- Index nodes from 1
- Children of node  $i$  will be  $2*i$  and  $2*i + 1$

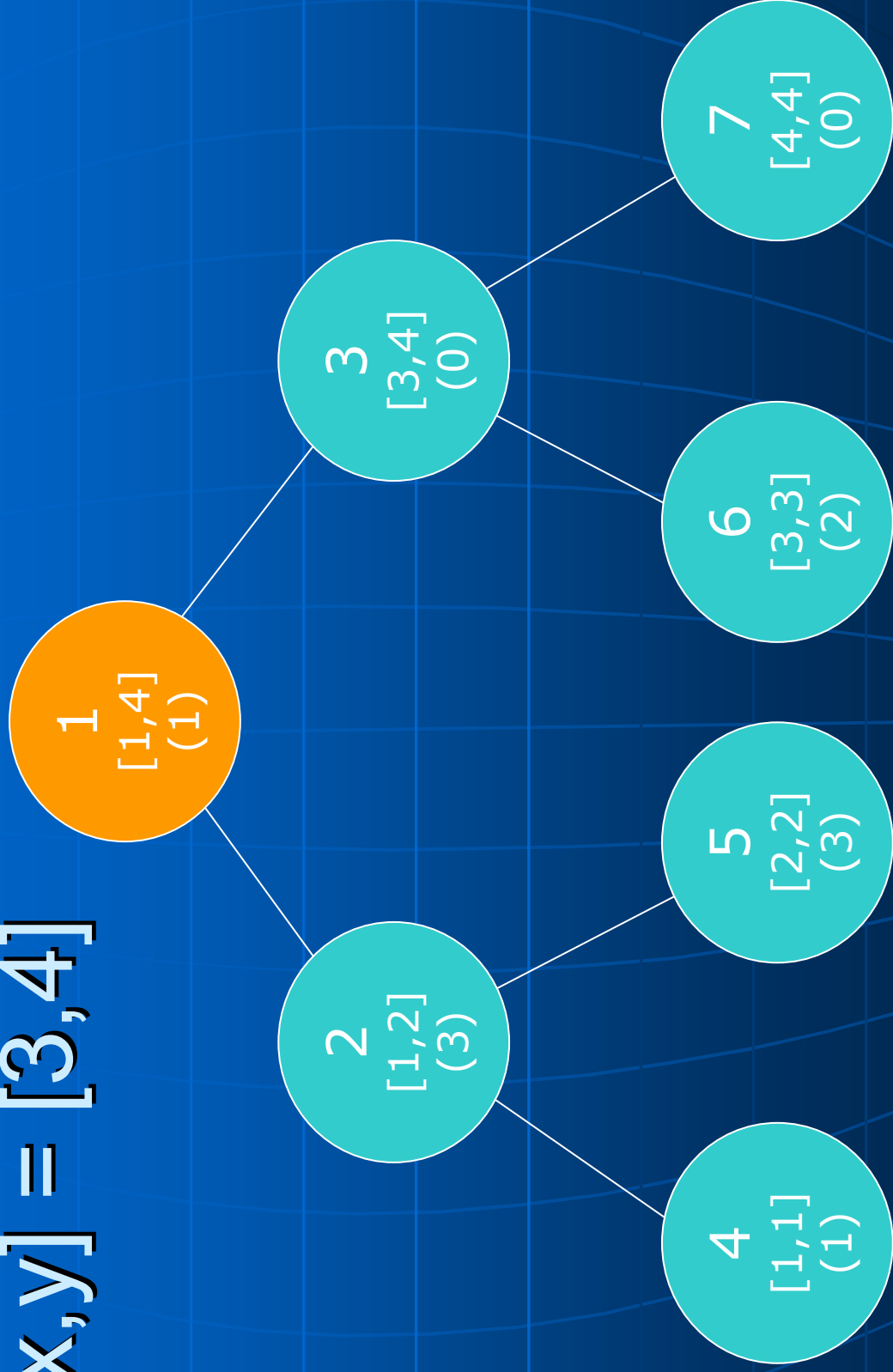
# Interval Trees: Updating a range





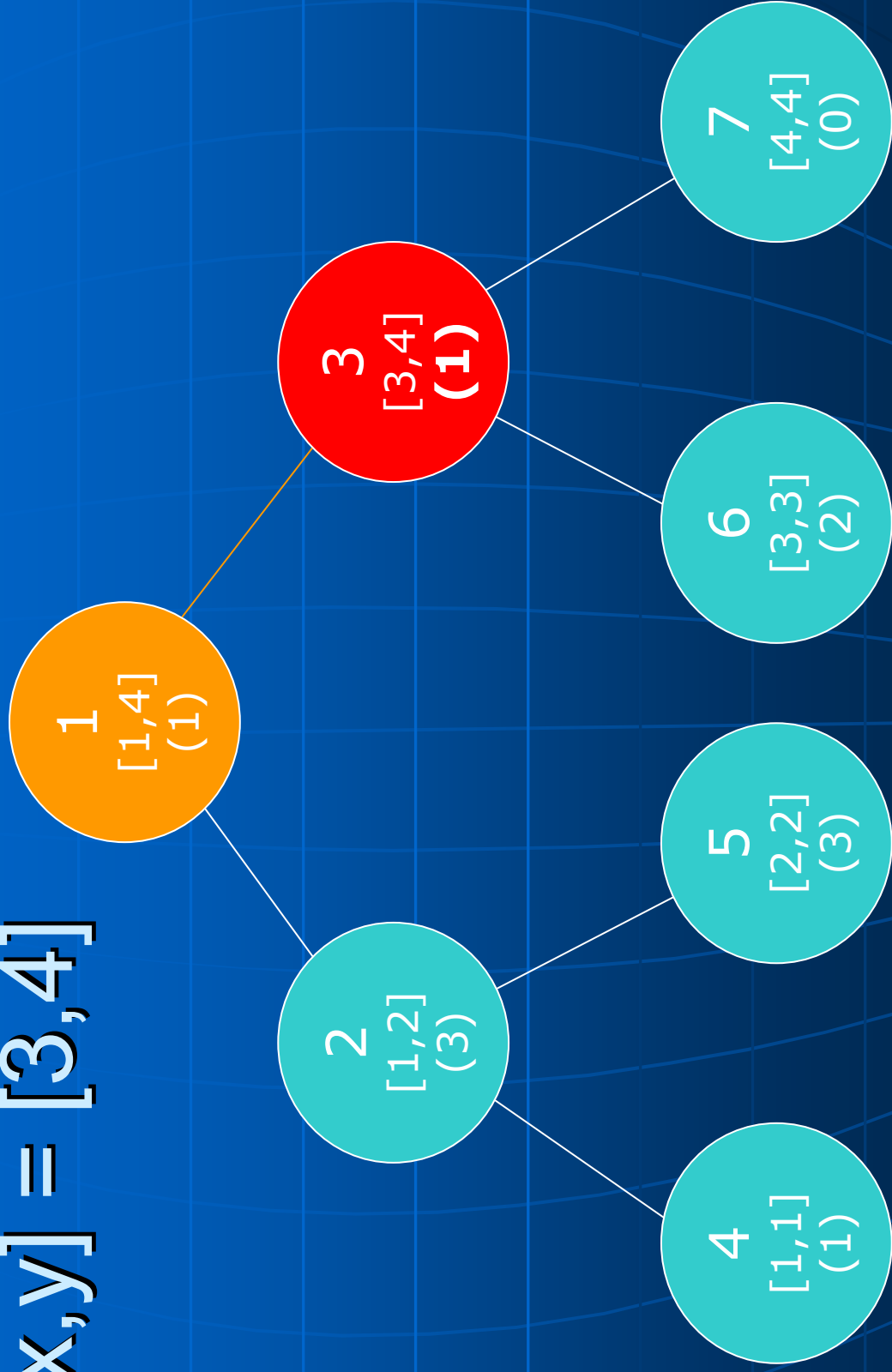
# Interval Trees: Updating a range

$[x,y] = [3,4]$



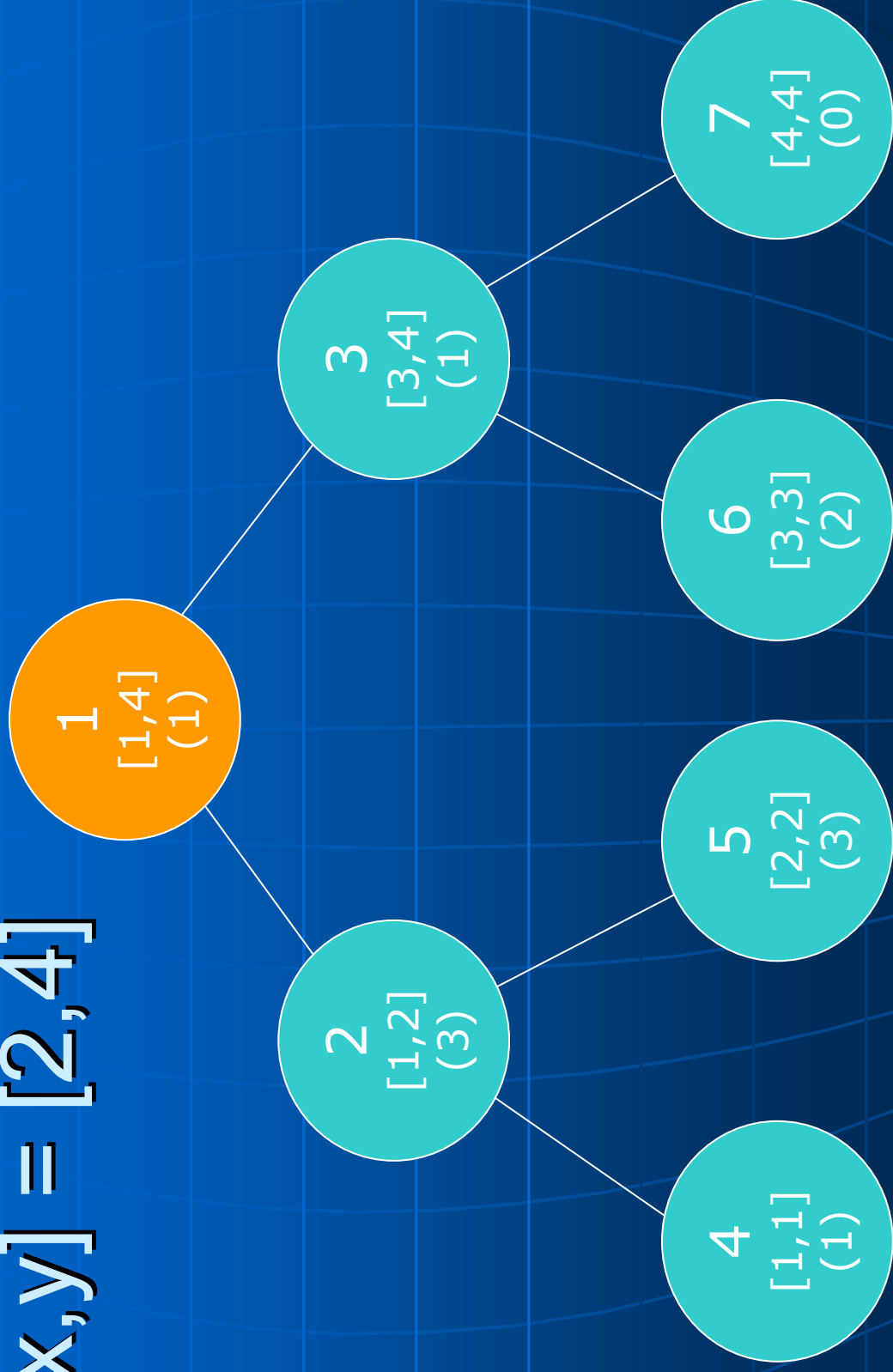
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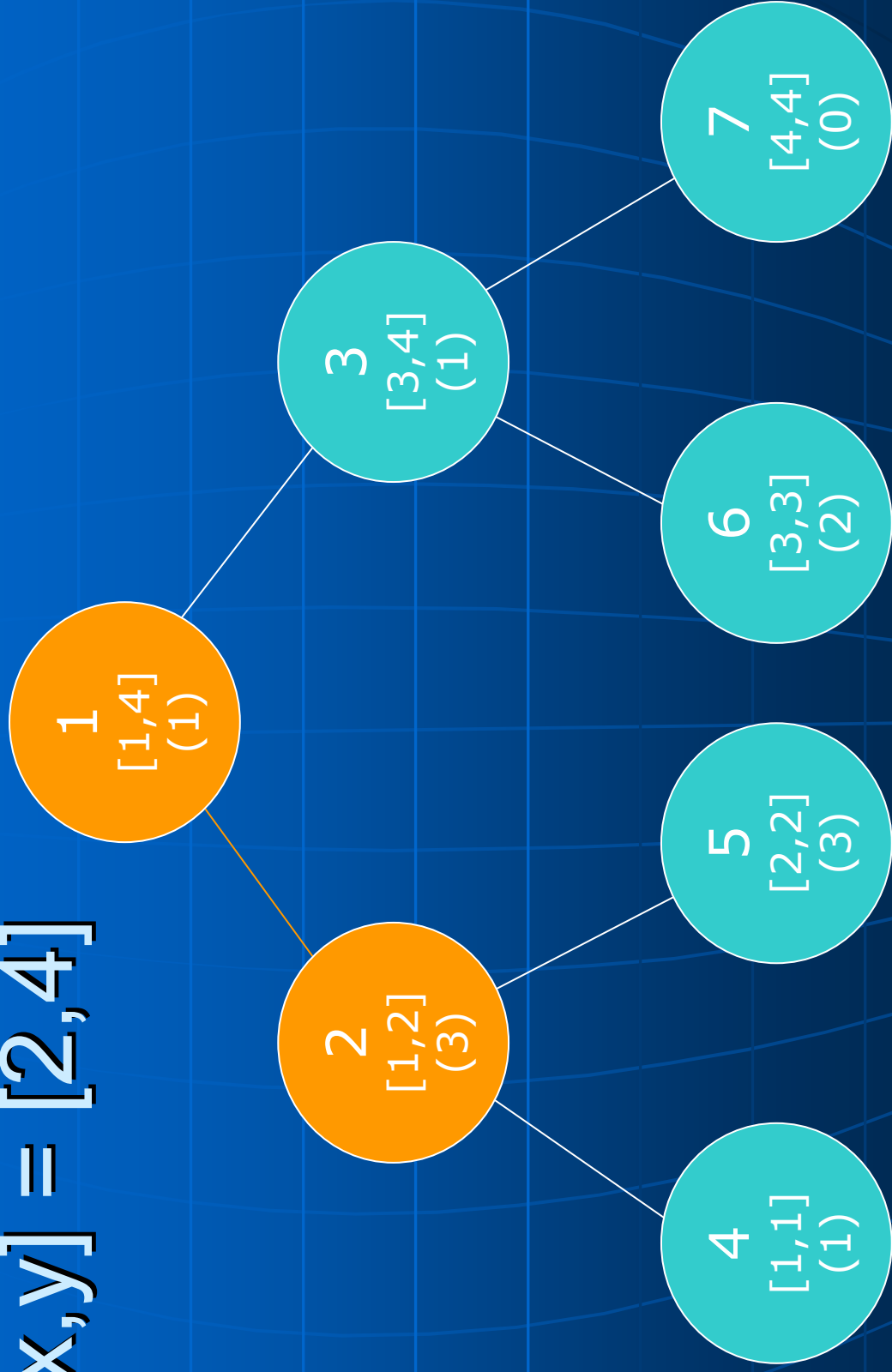
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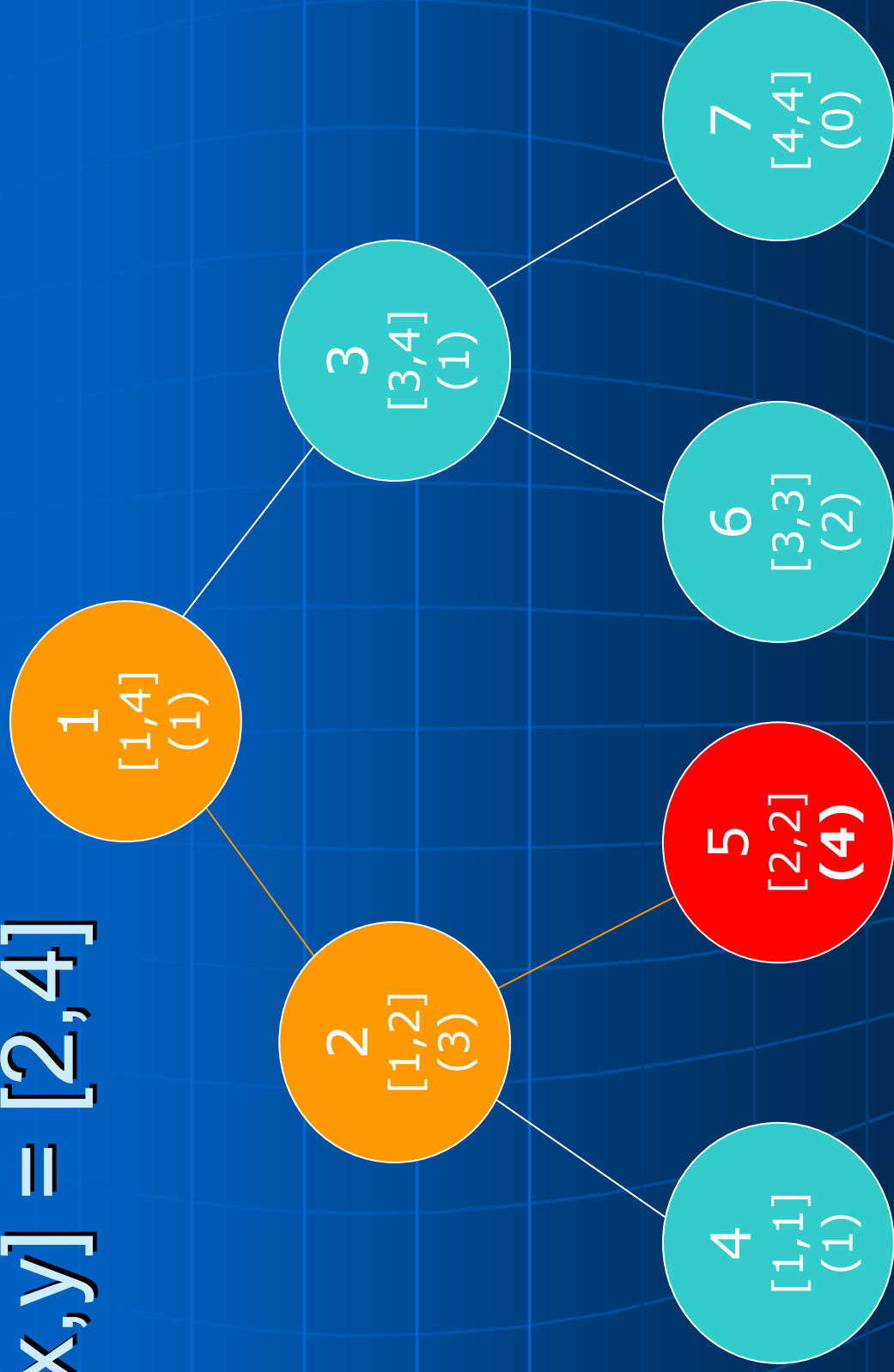
# Interval Trees: Updating a range

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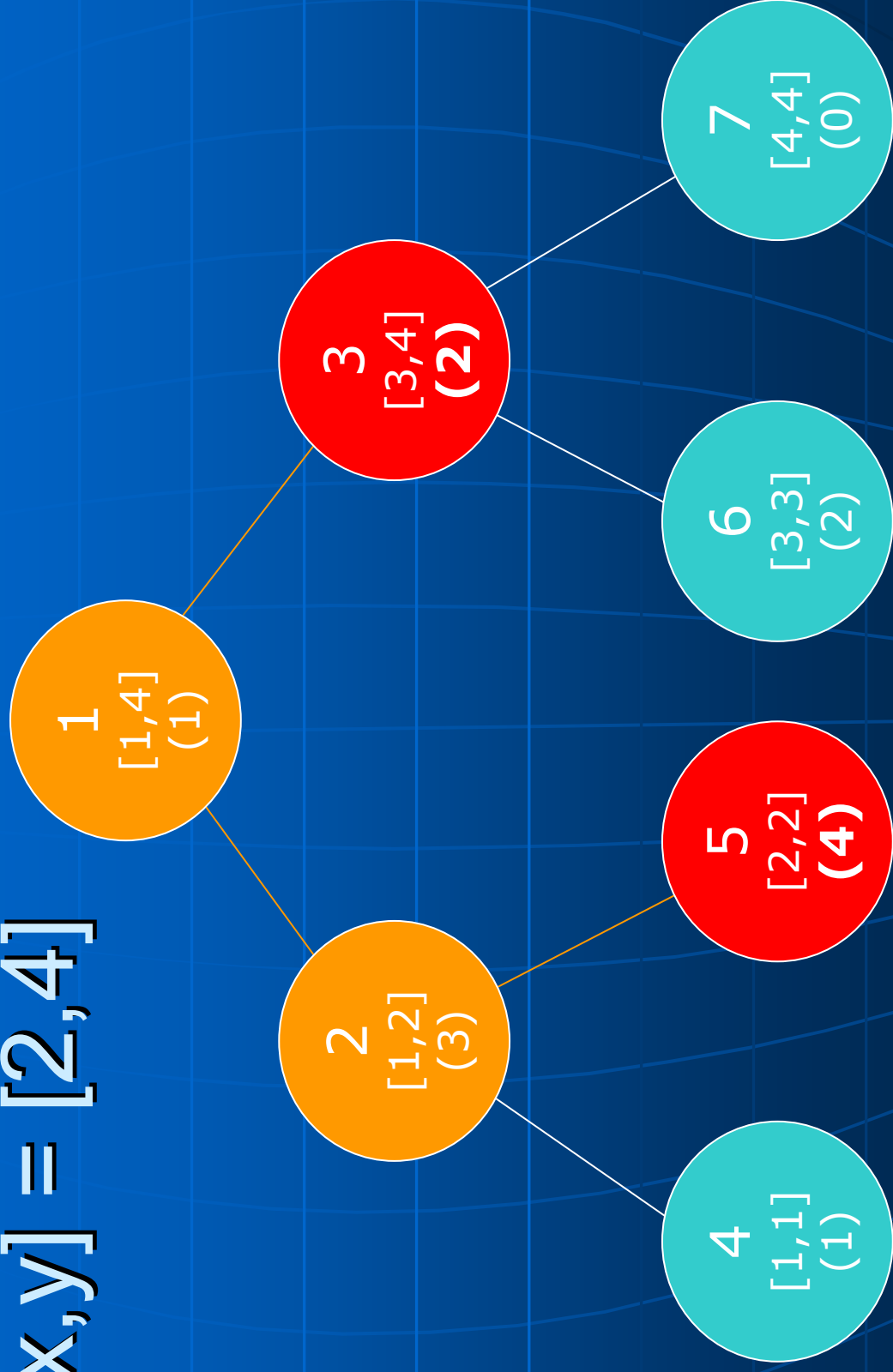
# Interval Trees: Updating a range

$[x,y] = [2,4]$



# Interval Trees: Updating a range

$[x,y] = [2,4]$



# Interval Trees: Updating Code

Updating a range  $[x,y]$ :

- Recurse into the tree, starting at root
- Suppose  $[A,B]$  is the current interval being considered:

    If  $[x,y]$  overlaps  $[A,B]$ :

        Increment count of node

    Else if  $[A,B]$  contains none of  $[x,y]$ :

        Stop

    Else:

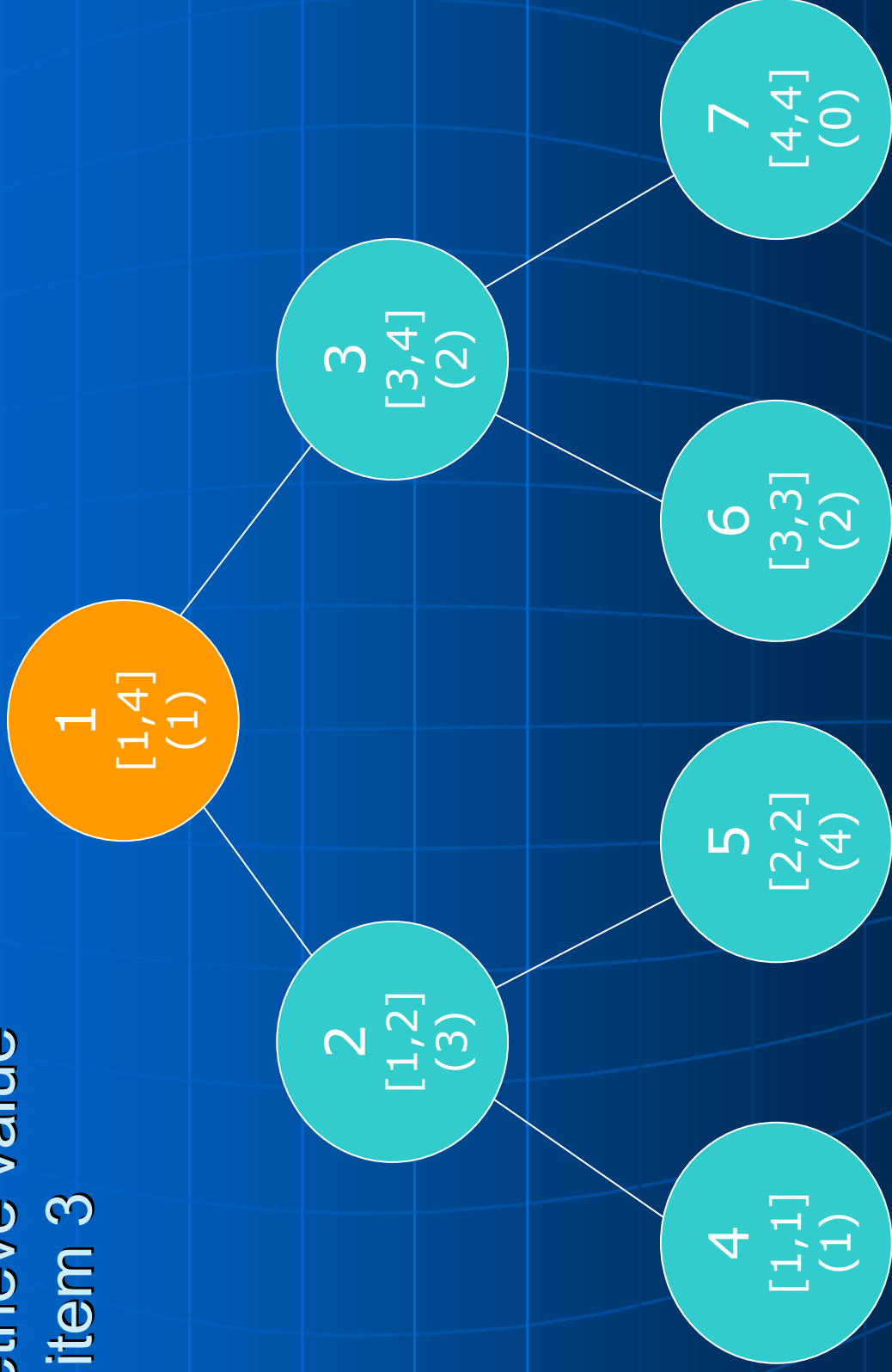
        Recurse into Left Child.

        Recurse into Right Child.

# Interval Trees: Query

Retrieve value  
of item 3

1

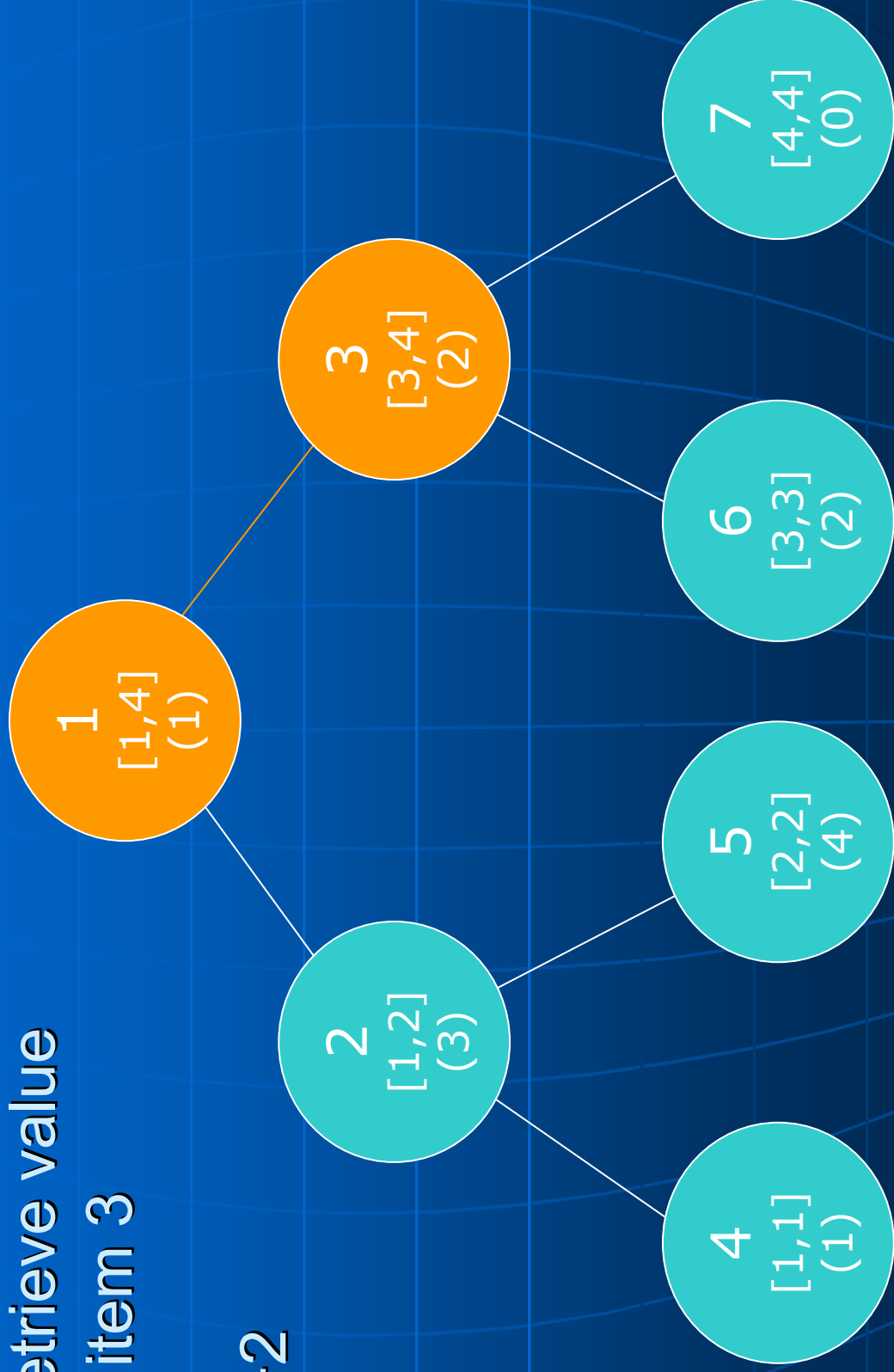




# Interval Trees: Query

Retrieve value  
of item 3

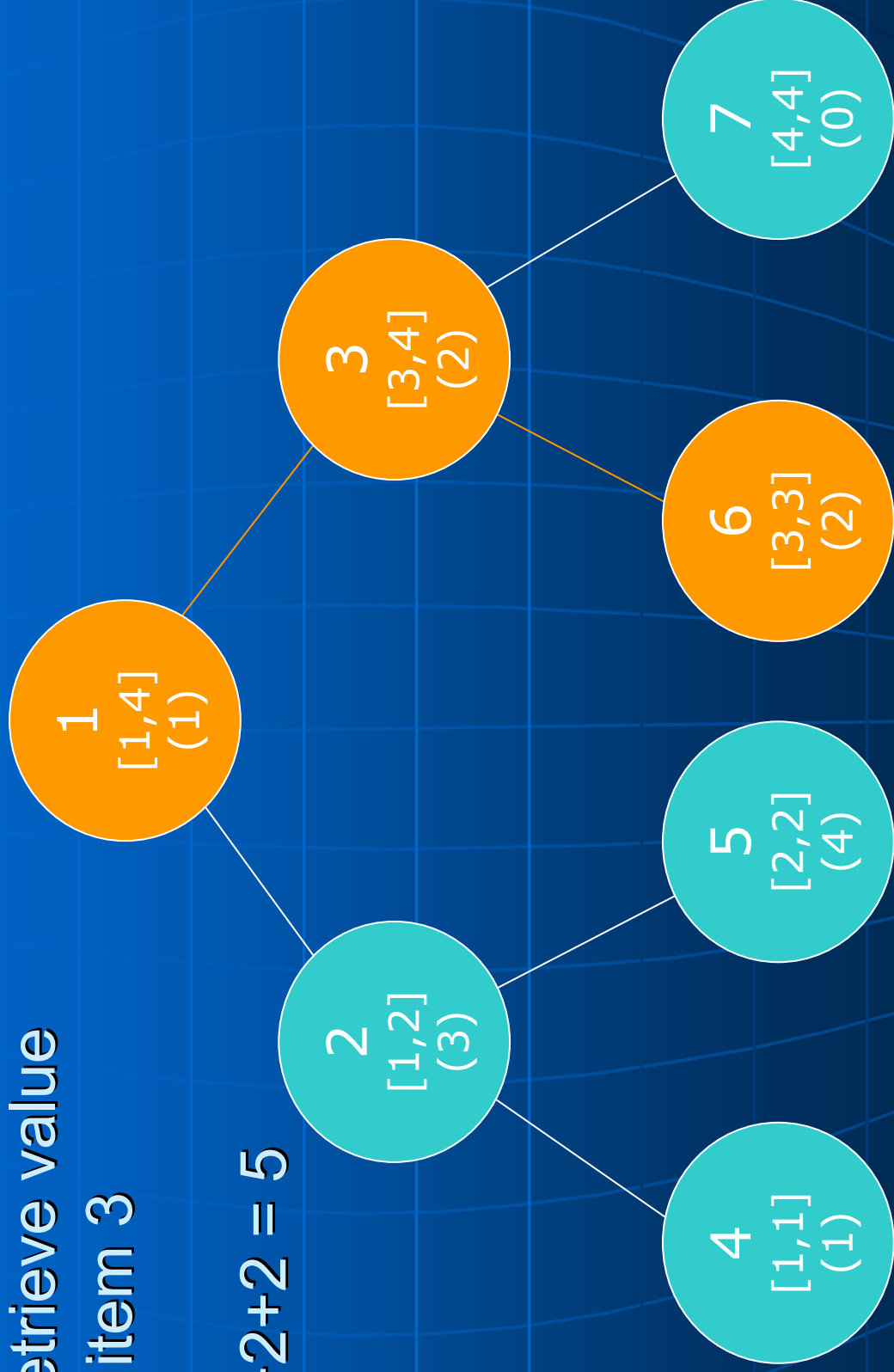
1+2



# Interval Trees: Query

Retrieve value  
of item 3

$$1+2+2 = 5$$



# Interval Trees: Query Code

Querying the number of item  $i$ :

- Let  $sum = 0$
- Recurse into tree.
- Considering range  $[A, B]$ :  
   $sum += \text{count of } [A, B]$
- If  $(A \neq B)$ :  
  Recurse into child containing  $i$   
  in its range.

# Binary Indexed Trees: Problem

Problem: Suppose we have a row of  $n$  numbers.

We want to:

- 1.) Increment value of the  $i$ th number
- 2.) Query the sum of the values in range  $[k, j]$

BIT implementation:

- Query can be done in  $O(\log n)$  for worst case
- Short code

Radix trees are equivalent in efficiency

# Binary Indexed Trees: Introduction

- Interested in counts over ranges of the form  $[1, i]$
- Uses the fact every number can be written as the sum of powers of 2.
- We will use this to represent ranges  $[1, i]$
- $13 = 8 + 4 + 1$
- $[1, 13] = [1, 8] + [9, 12] + [13, 13]$

# Binary Indexed Trees: Intervals

- We store interval counts of the form  $[i - 2^{r+1}, i]$  where  $r$  is the position of the last non-zero digit of  $i$  in binary form
- Example:
  - $13 = 1011$  in binary, position of last non-zero digit is 0.
  - $4 = 100$  in binary, position of last non-zero digit is 2.
- How it works will become clear when we explain the structure of the tree

# Binary Indexed Trees: Structure

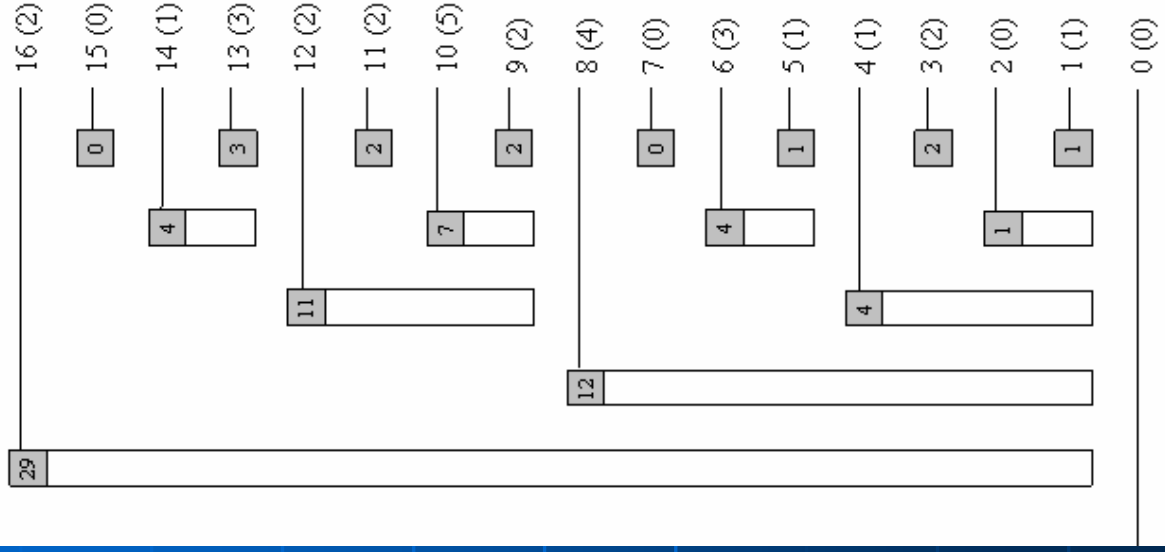
$f[i]$  = value at index  $i$

$c[i] = f[1] + f[2] + \dots + f[i]$

$tree[i] = \text{count of } [i-2^r + 1, i]$

$i$	$(i-2^{r+1}) \dots i$
1	1
2	1...2
3	3
4	1...4
5	5
6	5...6
7	7
8	1...8
9	9
10	9...10
11	11
12	9...12
13	13
14	13...14
15	15
16	1...16

# Binary Indexed Trees: Structure





# Binary Indexed Trees: Query

```
int read(int idx)
{
    int sum = 0;
    while (idx > 0)
    {
        sum += tree[idx];
        idx -= (idx & -idx);
    }
    return sum;
}
```

# Binary Indexed Trees: Update

```
void update(int idx ,int val)
{
    while (idx <= MaxVal)
    {
        tree[idx] += val;
        idx += (idx & -idx);
    }
}
```

# Questions

