



# Network Flow

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Maximum Flow

Minimum Cut



# What is Network Flow?

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- Graph Theory
- Flow
- Direction
- $A \rightarrow B$
- e.g. Water flowing through pipes



# Concepts and Definitions

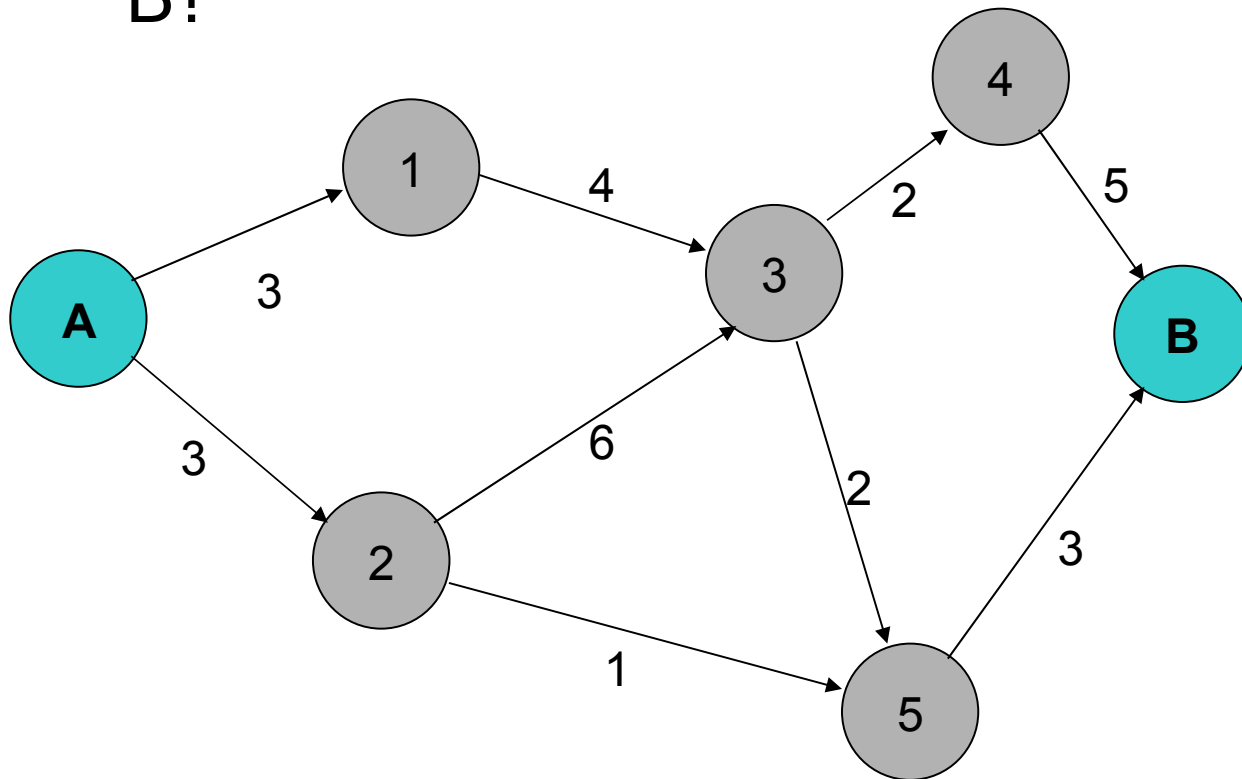
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- “Source” – Node in the graph emitting the “flow”
- “Sink” – Node in the graph consuming the “flow”
- “Capacity” – the maximum amount of “flow” that can pass through this edge/node.
- “Flow” – the amount of material passing through this node/edge/graph.

# Max Flow Problem

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- What is the maximum flow from A to B?



# Ford-Fulkerson

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- Simple

- Initialize all flow counts to 0.
- While there is an unsaturated path from source to sink:
  - Find the minimum capacity of all edges in that path
  - Increment the flow count of each edge in that path by the minimum
  - Increment the global flow count by that minimum
- The total amount of flow in the global flow count is then maximum

# Storage

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- For each edge store:
  - Capacity
  - Current flow
- Undirected
  - Depends on problem for optimal storage
  - Store 2 flows (positive and negative) or sometimes 2 capacities.
  - Depending on problem, positive may cancel out negative and vice versa.



# Finding the path

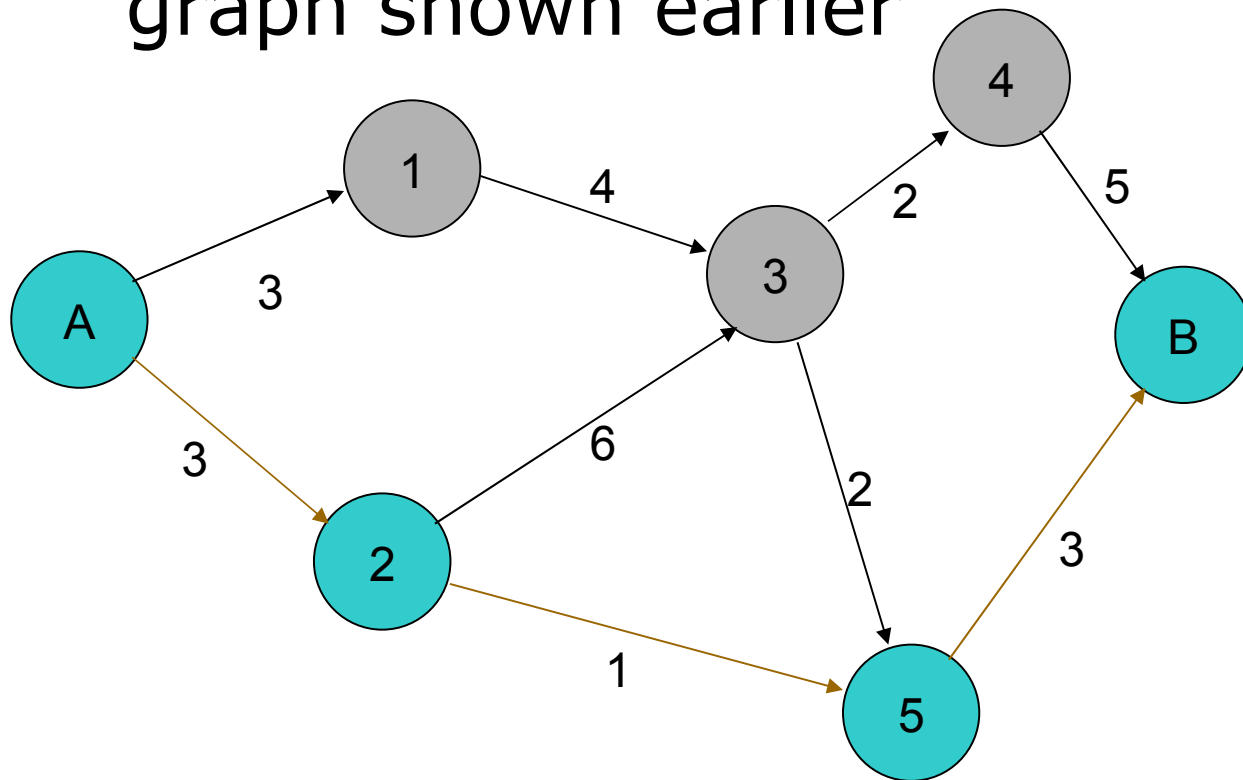
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- Method does not affect answer
- Affects running time
  
- Good choices (depending on problem)
  - Shortest path (BFS)
  - Path with maximum flow
  - DFS - simplest

# An example

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- Determining maximal flow in the graph shown earlier

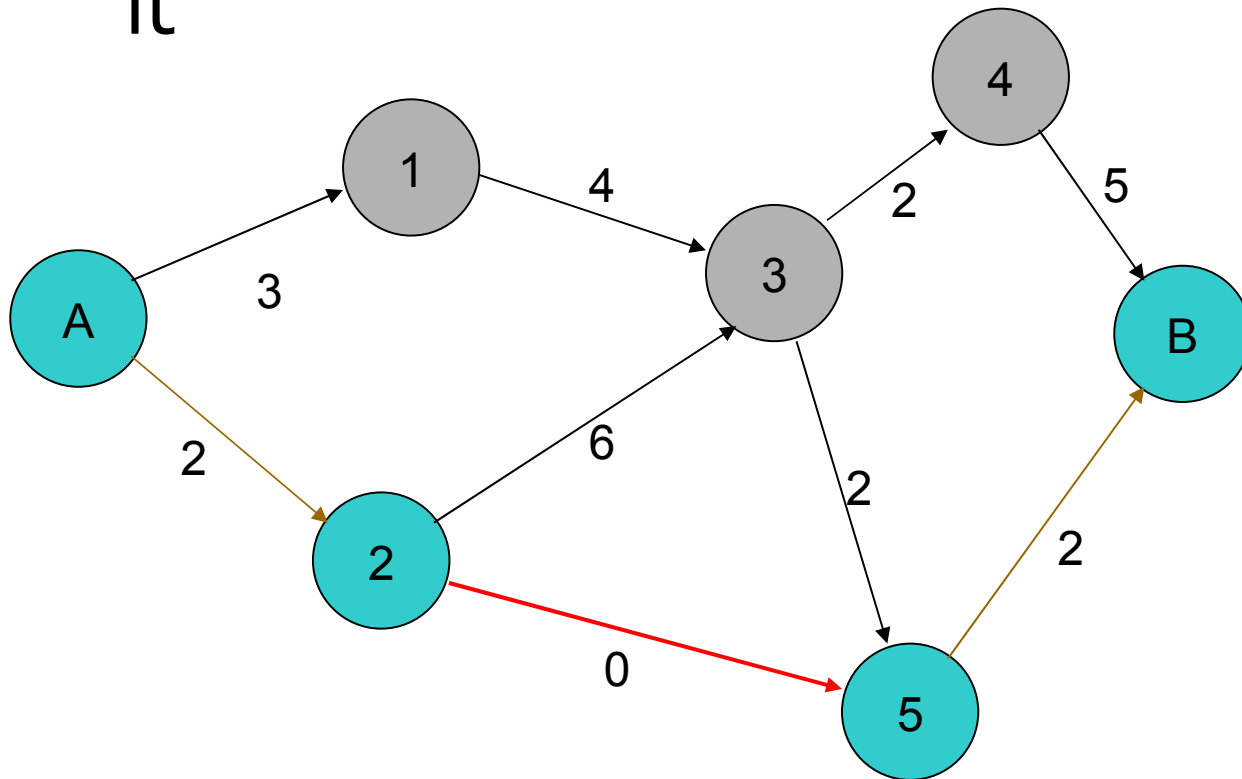




# An example

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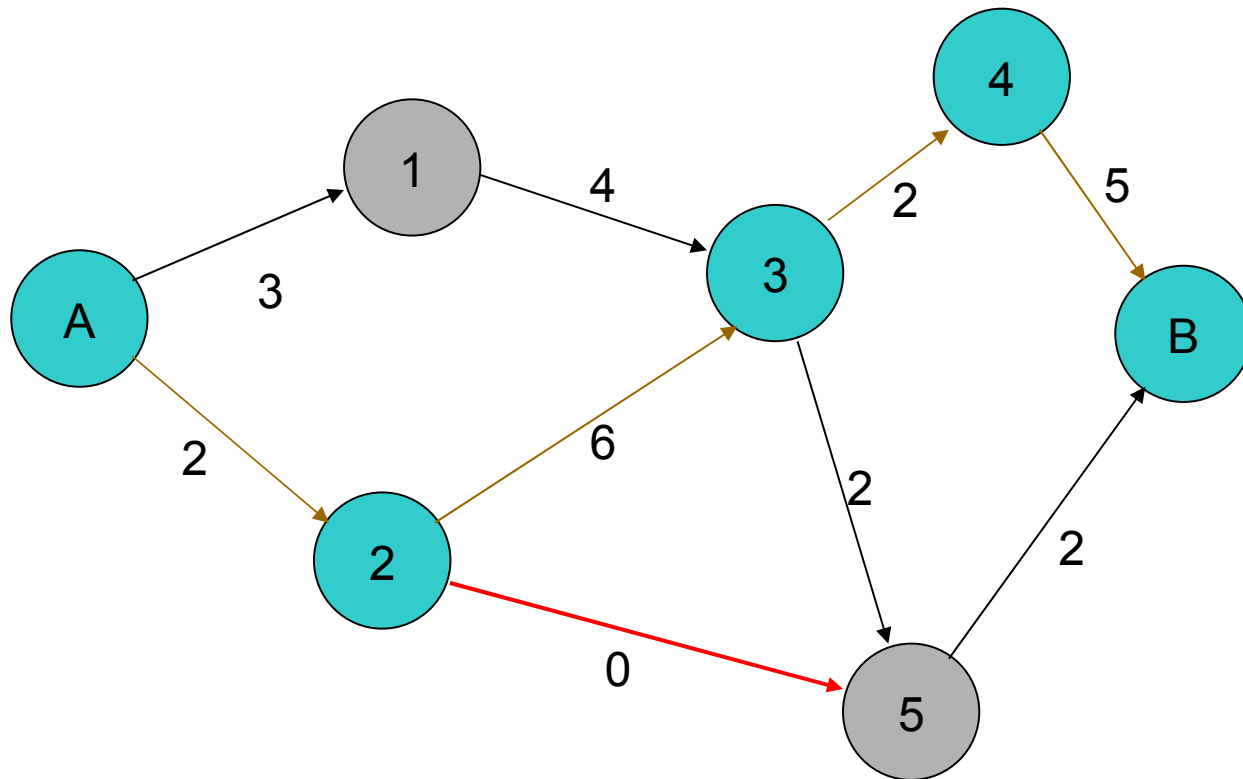
- Choose a path and put flow through it



# An example

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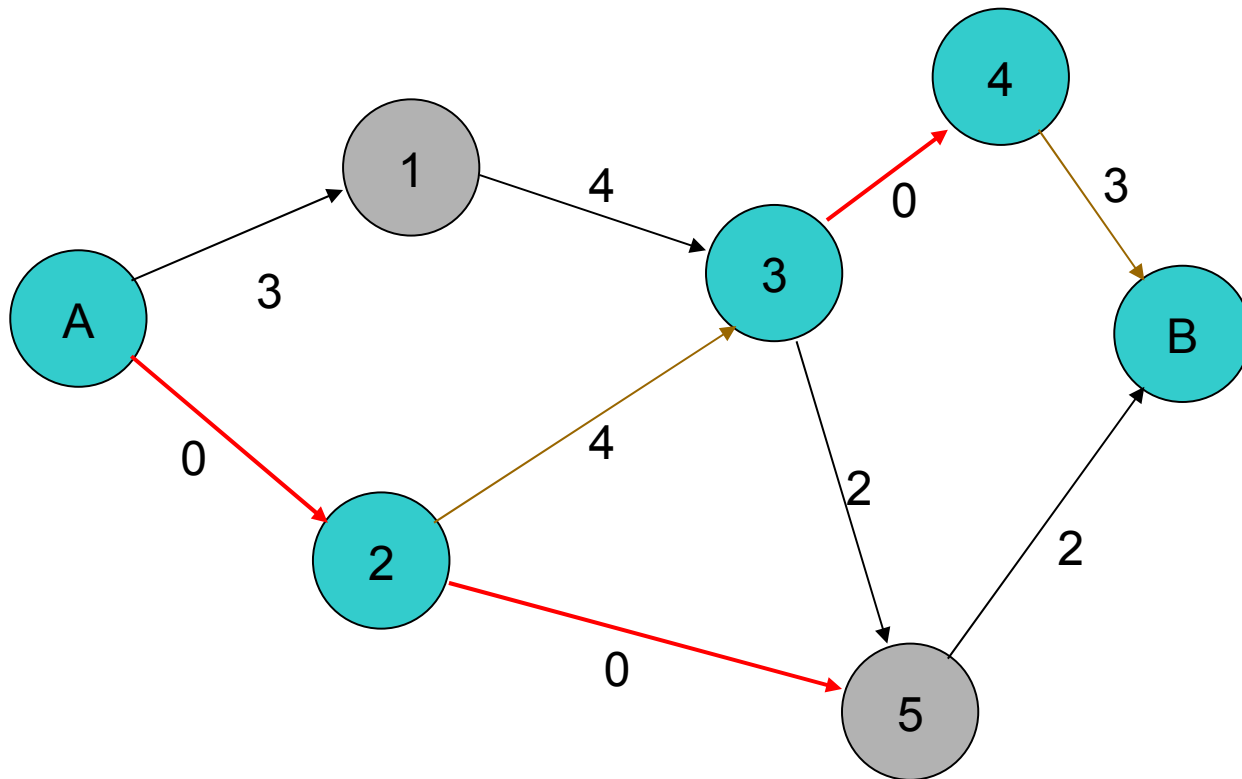
○ Current flow = 1



# An example

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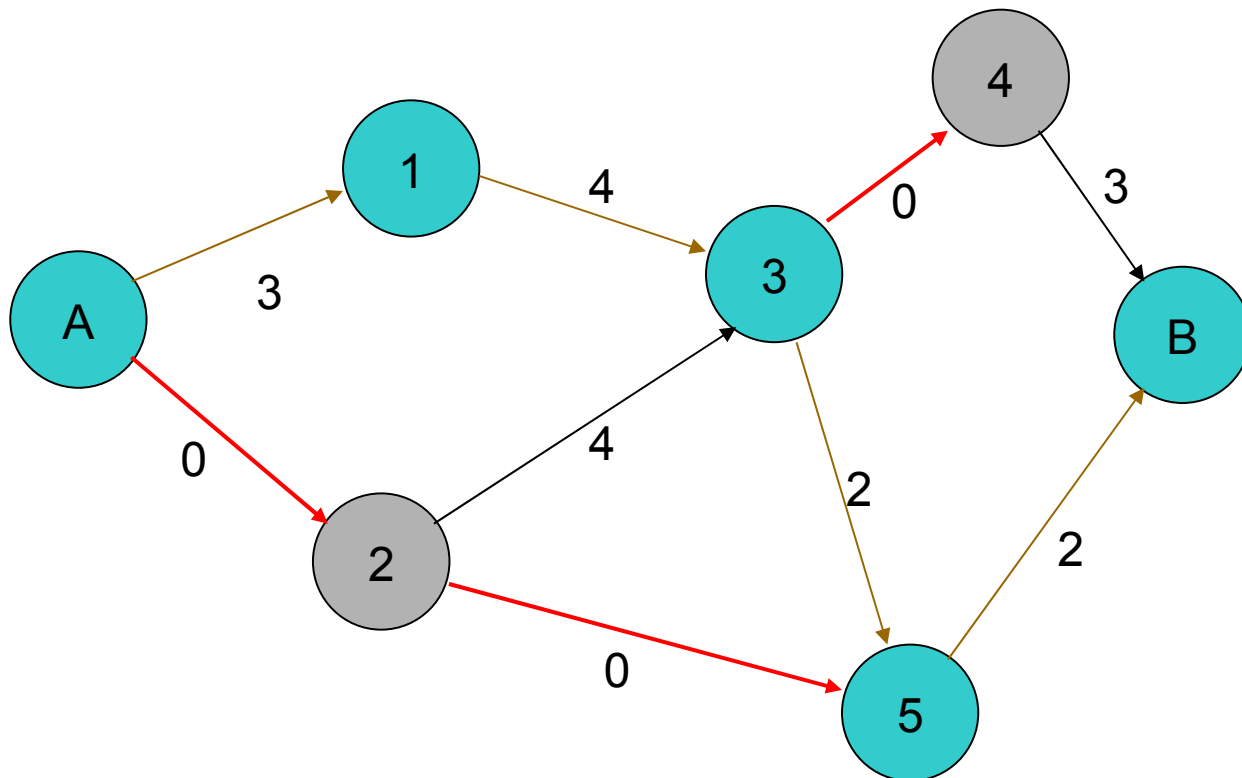
- Current flow = 3



# An example

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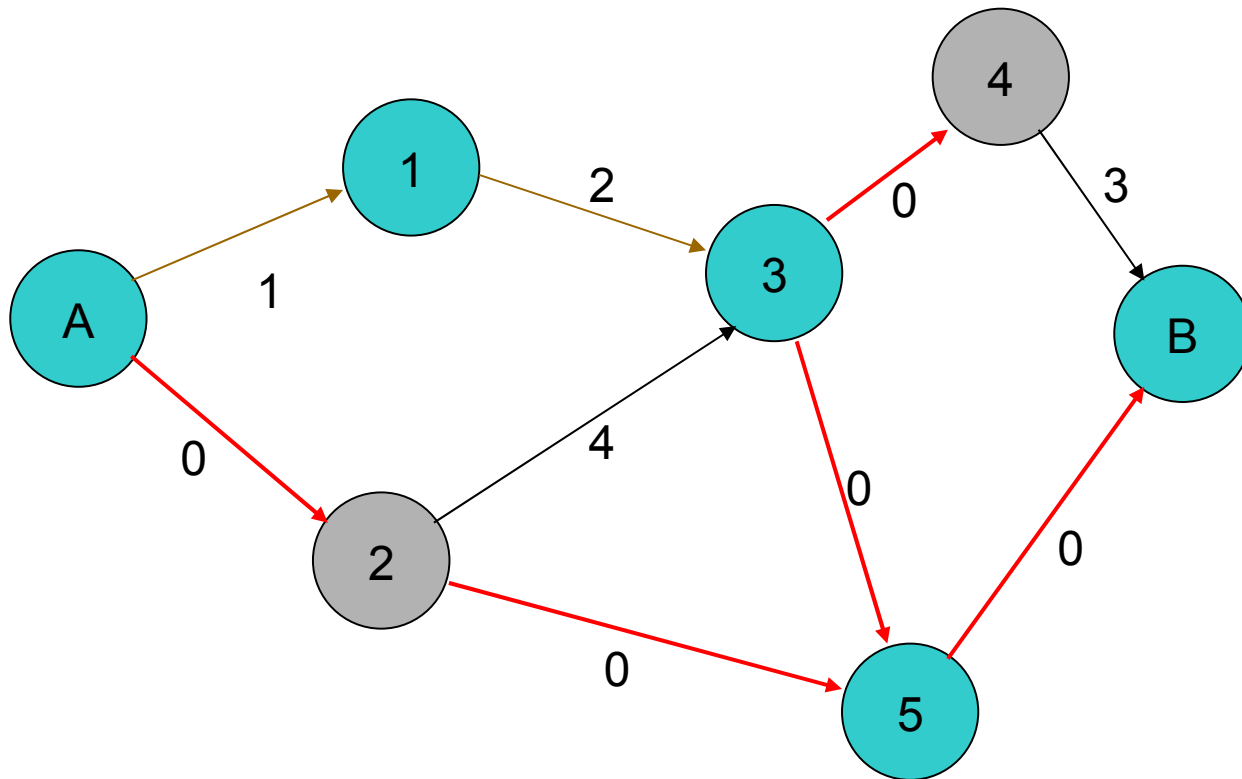
- Current flow = 3



# An example

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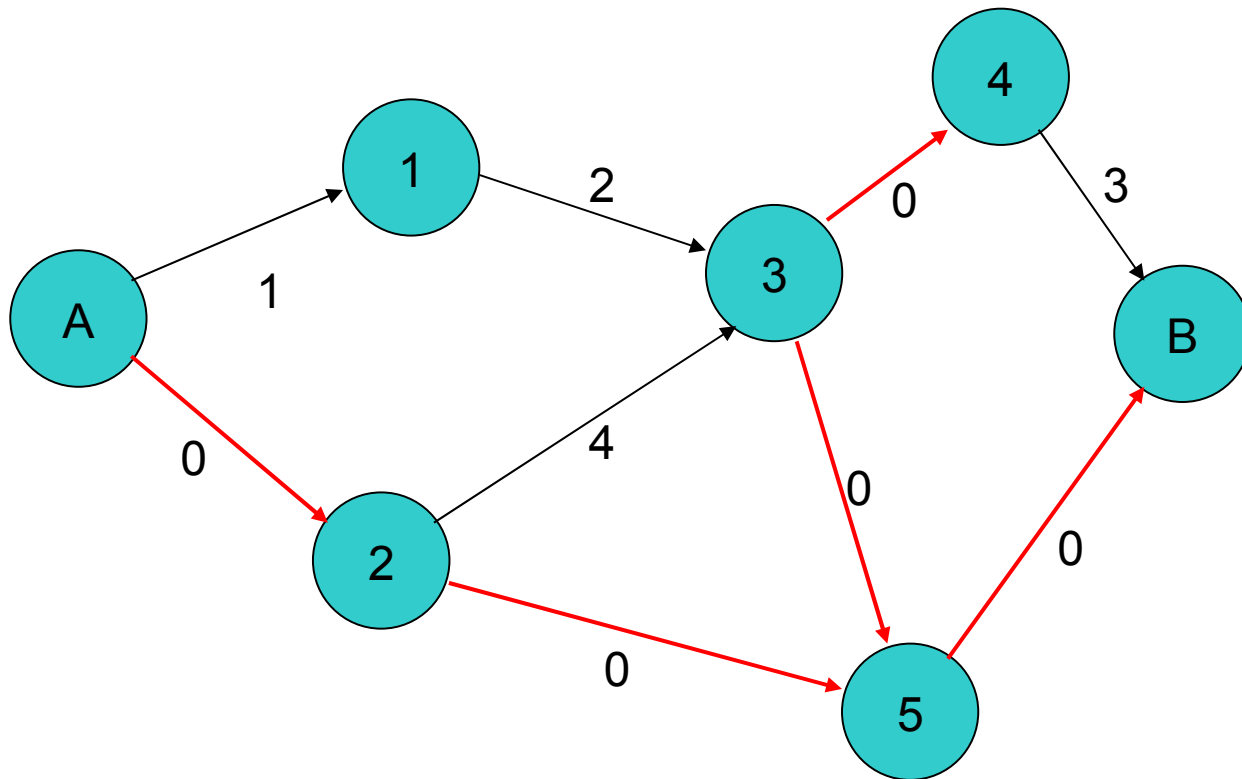
- Current flow = 5



# An example

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- Total flow = 5





# Minimum Cuts

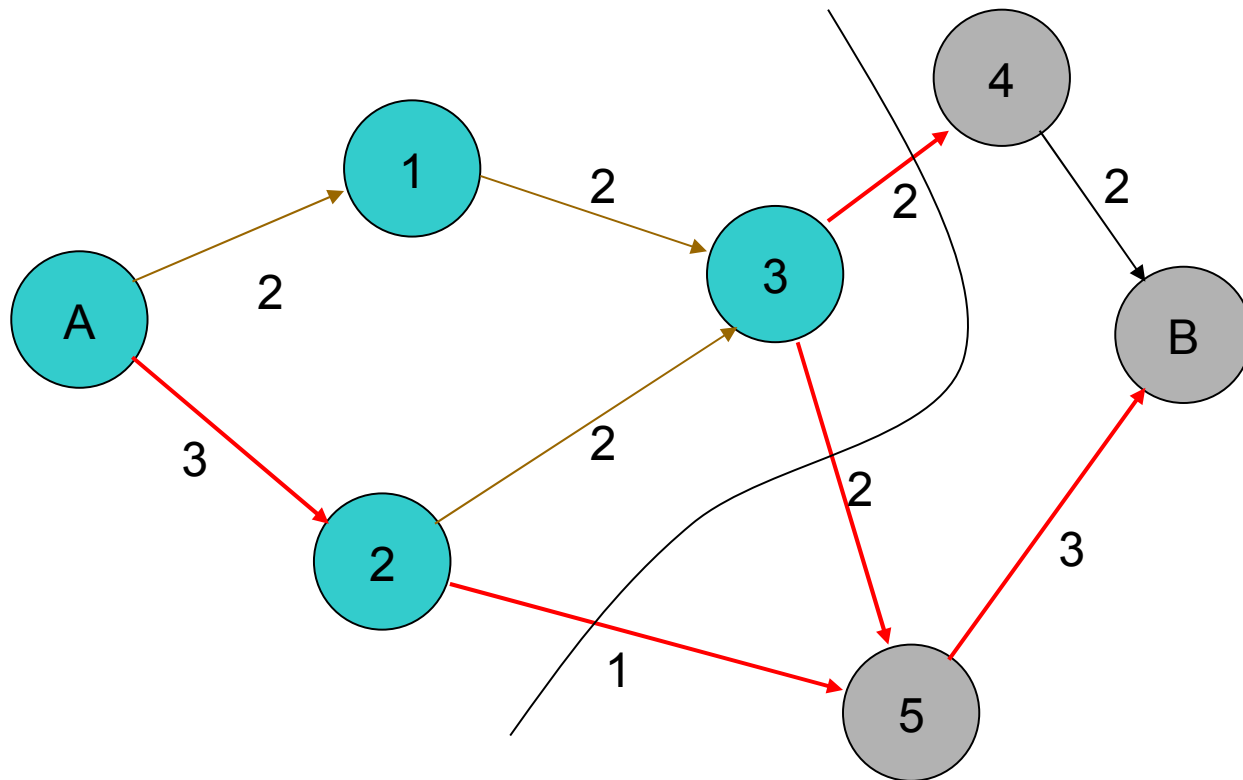
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- Maximum flow problems are tied in with Minimum Cut problems
- Minimum Cut
  - The minimum weighting of edges that separates two nodes (in this case source and sink)
- Sum of the weighting of cut edges = maximum flow through the network

# An example

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- Minimal cut =  $1 + 2 + 2 = 5$







## To find the minimum cut

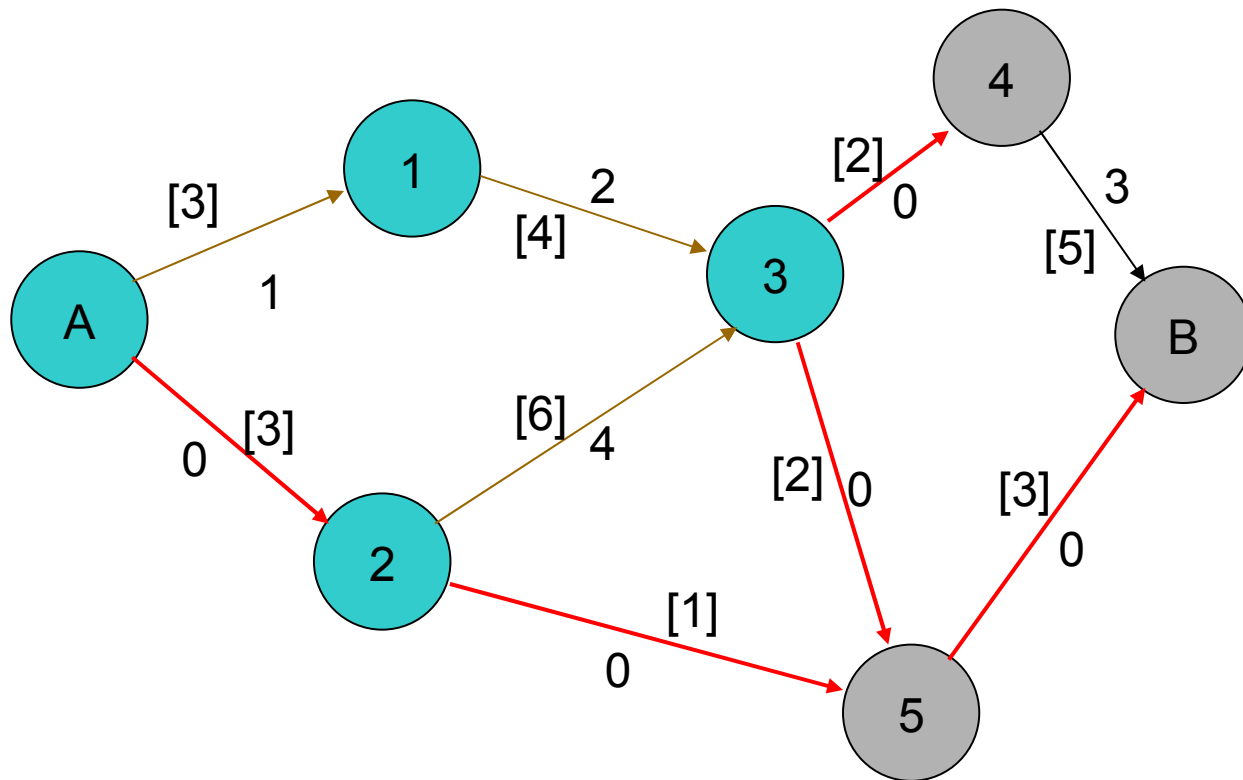
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- Create the maximum flow graph
- Select all nodes that can be reached from the source by unsaturated edges
- Cut all the edges that connect these nodes to the rest of the nodes in the graph
- This cut will be minimal

# An example

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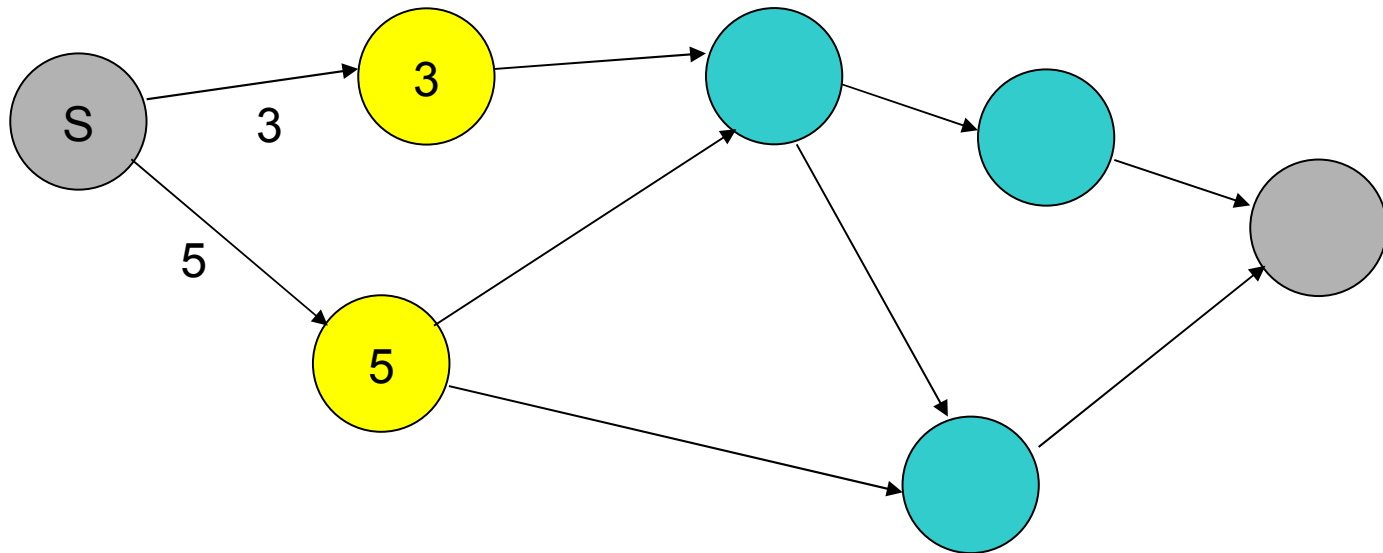
- To find a minimal cut



# Variations on Network Flow

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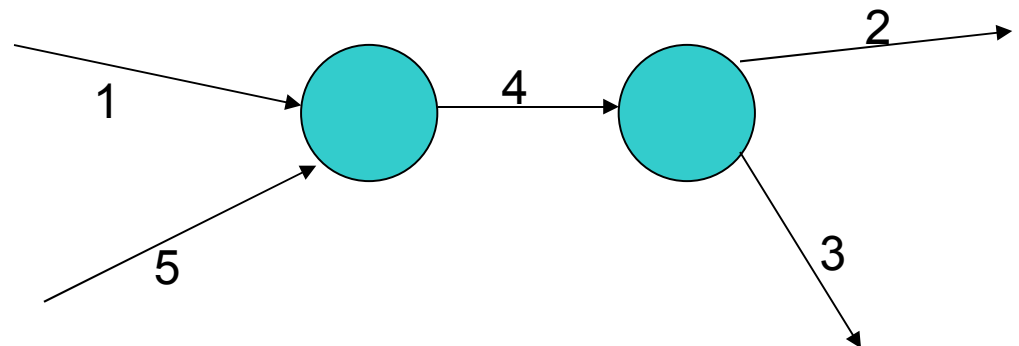
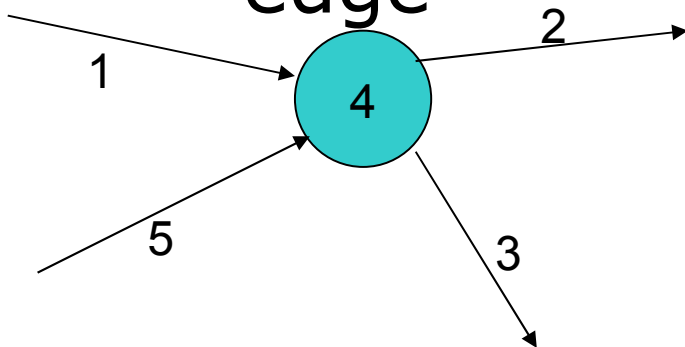
- Multiple sources & sinks – create a “supersource” or “supersink” which connects directly to each of these nodes, and has infinite capacity



# Variations on Network Flow

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- Node capacity – split the node into an “in” node, an “out” node and an edge



# Conclusion

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- Network flow problems
  - Graph Theory
  - Transporting some material from A to B
  - Along pathways (edges) that have capacities
  - Maximum flow is that maximum amount of material that can be transported from A to B.